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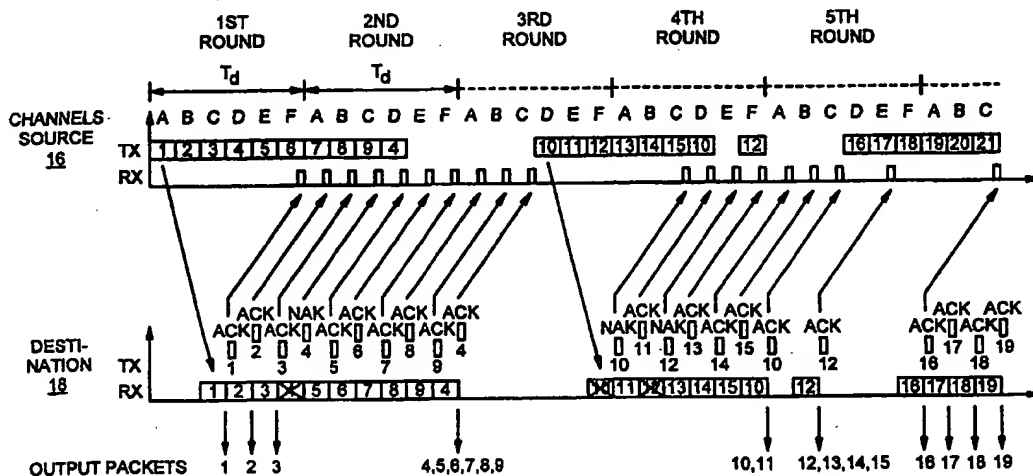
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(21) International Application Number: PCT/SE98/01415 (22) International Filing Date: 24 July 1998 (24.07.98) (30) Priority Data: 08/912,834          19 August 1997 (19.08.97)          US (71) Applicant: TELEFONAKTIEBOLAGET LM ERICSSON (publ) [SE/SE]; S-126 25 Stockholm (SE). (72) Inventor: HAARTSEN, Jacobus, Cornelis; Doddegras 29, NL-7623 DK Borne (NL). (74) Agent: ERICSSON MOBILE COMMUNICATIONS AB; Patent Unit, S-164 80 Stockholm (SE).		(81) Designated States: AL, AM, AT, AU, AZ, BA, BB, BG, BR, BY, CA, CH, CN, CU, CZ, DE, DK, EE, ES, FI, GB, GE, GH, GM, HU, ID, IL, IS, JP, KE, KG, KP, KR, KZ, LC, LK, LR, LS, LT, LU, LV, MD, MG, MK, MN, MW, MX, NO, NZ, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TR, TT, UA, UG, UZ, VN, YU, ZW, ARIPO patent (GH, GM, KE, LS, MW, SD, SZ, UG, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GW, ML, MR, NE, SN, TD, TG).  Published With international search report.	

(54) Title: A MULTI-CHANNEL AUTOMATIC RETRANSMISSION QUERY (ARQ) METHOD



(57) Abstract

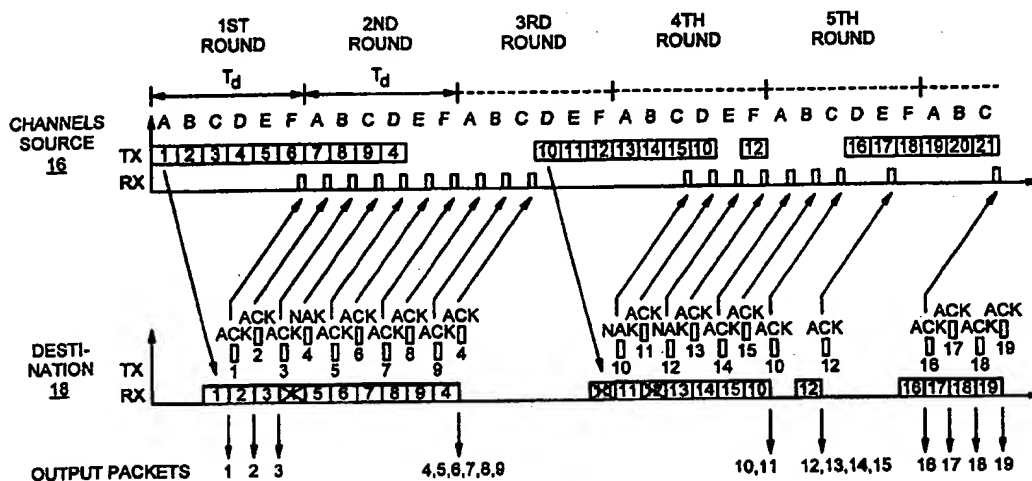
A multi-channel Automatic Retransmission Query (ARQ) method transmits data packets from a source to a destination over a communication link that is subdivided into a number of channels. A network using the multi-channel ARQ method sequentially multiplexes the data packets at the source and transmits them over corresponding channels. The network applies a stop-and-wait ARQ method on each one of the channels and determines whether the destination has positively acknowledged a previously transmitted data packet. If not, the network retransmits only the data packets that are not positively acknowledged.



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A multi-channel Automatic Retransmission Query (ARQ) method transmits data packets from a source to a destination over a communication link that is subdivided into a number of channels. A network using the multi-channel ARQ method sequentially multiplexes the data packets at the source and transmits them over corresponding channels. The network applies a stop-and-wait ARQ method on each one of the channels and determines whether the destination has positively acknowledged a previously transmitted data packet. If not, the network retransmits only the data packets that are not positively acknowledged.

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## A MULTI-CHANNEL AUTOMATIC RETRANSMISSION QUERY (ARQ) METHOD

### Field of Invention

This invention relates to the field of data  
5 communications, more particularly, to improving data  
throughput in a data communication network.

### Background

Today's information age has experienced a tre-  
10 mendous growth in the volume of data communicated over  
various networks. For example, enormous amounts of data are  
transferred every day between computers connected by data  
networks like intra-nets, internet, local area nets, wide  
area nets, etc. The speed by which these data are  
15 communicated, i.e., data throughput, is an important  
consideration in a data network. Of course, the network's  
data throughput is proportional to the amount of  
communication resources allocated for communicating data,  
with more allocated resources resulting in a higher data  
20 throughput. Since the communication resources are scarce,  
many networks share these resources among many users.

Another important consideration is a network's  
overall communication error rate, with some networks  
requiring residual error rates of less than  $10^{-16}$ . Some  
25 communication networks use packet-switched communication  
links in which the user information is packetized, and the  
packets are individually sent as data burst from a source to  
a destination. In packet-switched networks, the data

throughput is expressed in terms of the number of data packets per unit of time. Since the links are not perfect, especially when wireless links are used, the error rate requirement is not easily met unless some type of error  
5 correction mechanism is introduced, for correcting errors while data is being communicated over the network.

Error correction schemes can be divided into two classes: a forward-error-correction (FEC) method, and a retransmission method. In the FEC method, extra coding bits  
10 are added at the source to increase the redundancy in the information. This extra redundancy will enable the destination to correct errors. In the retransmission method, extra check bits are added to the data packets at the source so that the destination can check the data  
15 packets for errors. If the destination detects an error, it will automatically request for a retransmission of the data packet. This type of retransmission method is known as Automatic Retransmit Query (ARQ) method.

The data traffic over a switched packet network is  
20 carried out in data bursts. Communication errors often arrive in bursts as well. Therefore, it is more efficient to apply a ARQ method, which requires occasional retransmissions, than to apply FEC coding, which adds a constant overhead irrespective of whether there are  
25 communication errors or not. A slight amount of FEC coding is always advantageous to avoid retransmissions of data packets that have very few errors. Accordingly, the use of an ARQ method, preferably, with little FEC, is necessary in

order to produce a low residual error rate in a high data throughput network.

A number of ARQ methods is known. In general, the ARQ methods may be divided into a numbered ARQ method and an  
5 unnumbered ARQ method. In the numbered ARQ method, each data packet is assigned a packet number that allows the destination to specify which data packets have arrived successfully and which have arrived erroneously. In this way, the data packets do not have to arrive at the  
10 destination in the order of transmission from the source. As is well known, the packet-switched networks transfer consecutive packets along different paths having different delays. Therefore, the data packets may arrive out of sequence. By using the packet numbers of the numbered ARQ  
15 method, the data packets can be reordered at the destination in proper sequence.

In the unnumbered ARQ method, packets do not carry a packet number. Therefore, the unnumbered ARQ method requires less overhead and, thus, is less complex to  
20 implement. In order to avoid out-of-order arrival of data packets, however, the data packets must be communicated in an orderly manner. A well-known unnumbered ARQ method is a stop-and-wait ARQ method in which the source stops and waits until a transmitted data packet is acknowledged. For each  
25 data packet, a positive acknowledgement (ACK) must be received from the destination, before a subsequent data packet can be transmitted from the source. If a negative acknowledgement (NAK) is received from the destination,

however, the source retransmits the same data packet again. If no acknowledgement is received, the source will automatically retransmit the same data packet after a time-out period.

5           The delay between the data packet transmission and the arrival of the acknowledgement information is known as a round-trip delay. The round-trip delay determines the data throughput of the network. The longer the round-trip delay, the longer the source has to wait before it can transmit a  
10 new data packet or retransmit a NAKed data packet. Accordingly, the data throughput on the communication link is inversely proportional to the round-trip delay between the source and destination. Because during the wait  
15 periods, the source is idle and no transmission takes place, in order to maintain a high data throughput, the stop-and-wait ARQ methods are generally used in networks with short round-trip delays.

FIG. 1 shows an exemplary timing diagram of a data packet transfer that uses a conventional stop-and-wait ARQ  
20 method. As shown, a sequence of five (5) data packets are transmitted from a source 12 to a destination 14. The data packet are shown by numbers 1 to 5 for illustration purpose only. In practice, however, no number is assigned to the data packets. The round trip delay between a packet and its  
25 response is shown by  $T_d$ . If a data packet is not received successfully, the destination 14 transmits a NAK to the source. For example, an erroneously transmitted 3rd data packet, shown by number 3 in FIG. 1, is retransmitted from

retransmitted again. Although less so than ordinary go-back-N ARQ method, the cumulative method does not prevent the retransmission of successfully received data packets. The buffer lengths in both of these methods depend on the  
5 round-trip delay. As in the stop-and-wait ARQ method, the data throughput decreases, when the round-trip delay increases, but only in an error prone environment. Under error free conditions, optimal throughput is obtained.

The most efficient ARQ scheme, both with respect  
10 to throughput per user and shared resource efficiency is a selective repeat ARQ method. According to this method, the destination specifies in the acknowledgement exactly which data packet has to be retransmitted. In this way, only failed packets are retransmitted, and no packet transmission  
15 is delayed due to a retransmission. The packet numbers are used both for retransmission requests and for the ordering of out-of-sequence data packets. The data throughput is independent of the round-trip delay, and can reach the theoretical limit of what can be achieved with a  
20 retransmission method, that is a data throughput of  $1 - FER$ , where FER is Frame Error Rate or Packet Error Rate. However, the selective repeat ARQ method requires extensive storage capability and packet handling at both the source and the destination.

25 Some communication networks transfer data packets in parallel. These networks subdivide the communication link into a number of subchannels for parallel transmission of data packets. In order to guarantee proper ordering of



data packets arriving from various subchannels, most conventional networks employ the numbered ARQ method. One conventional network, however, uses the unnumbered ARQ method for transfer of data packets in parallel. This network, which is used by ARPANET and is described by D. Bertsekas and R. Gallager, "Data Networks," 2nd Edition, Prentice Hall International, Inc., London, 1992, applies independent stop-and-wait ARQ methods on eight (virtual) channels in parallel. The data packets are multiplexed over the channels according to availability to transmit a new data packet on the channels. Each data packet contains a virtual channel number identifying the channel on which the data packet is transmitted. As a consequence, the data packets can be transmitted out of order, because the stop-and-wait ARQ method on the separate channels are independent of each other. In other words, retransmissions on one channel will not affect the other channels, but will result in data packets arriving at the destination out of order. As a result, packet numbers at higher protocol layers as required to provide enabling instructions to the destination to properly, reorder the data packets. This method, therefore, requires complicated data-packet reordering overhead, which must be implemented in higher protocol layers.

Therefore, there exists a need for a data communication network with high data throughput which is simple to implement and does not retransmit successfully received data packets unnecessarily.

### Summary

Briefly, the present invention is exemplified in a communication network that uses a multi-channel automatic retransmission query (ARQ) for communicating data packets from a source to a destination over a communication link that is subdivided into a number of channels. The network sequentially multiplexes the data packets at the source over corresponding channels during multiplexing rounds. The invention then applies a stop-and wait ARQ method on each one of the channels and determines whether the destination has positively acknowledged a previously transmitted data packet. If not, the network retransmits only the data packets that are not positively acknowledged, after being transmitted during a previous multiplexing round.

At the source, the network of the invention multiplexes the data packets during successive multiplexing rounds, when the data packets are sequentially transmitted on all of the channels. The network halts the multiplexing of new data packets at the source during a subsequent multiplexing round until the destination positively acknowledges successful reception of a data packet that was transmitted over a corresponding channel during the previous multiplexing round. The network only retransmits previously transmitted data packets that are not positively acknowledged. Also, the source may retransmit the data packets if no acknowledgement is received from the destination after a predefined time-out period. At the destination, the network sequentially de-multiplexes the

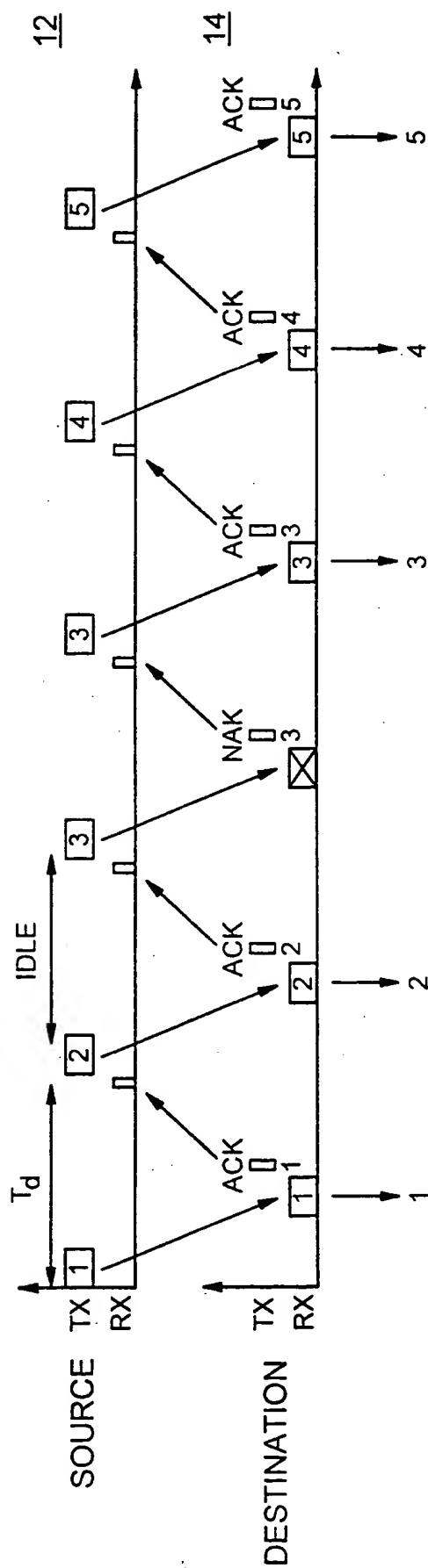
FIG. 4 is a block diagram implementing a multi-channel ARQ method according to another embodiment of the present invention.

5 FIG. 5 (a) and 5(b) are diagrams showing two exemplary operation of the multi-channel ARQ method of the embodiment of FIG. 4.

FIG. 6 is a block diagram implementing multi-channel ARQ  
10 method according to yet another embodiment of the present invention.

#### Detailed Description

Referring to FIG. 2, a multi-channel ARQ method  
15 according to the present invention subdivides each round trip delay into a number of time slots. As shown, the roundtrip delay  $T_d$  is divided into six time slots designated as ARQ Channels A-F. In a Time Division Multiplex (TDM) manner, each time slot is used as a virtual channel on a  
20 communication link assigned to a user. These channels are used for parallel transfer of data packets between a source 16 and a destination 18. At the source 16, the data packets are sequentially multiplexed over the ARQ Channels A-F and transmitted to the destination 18. The multi-channel ARQ  
25 method of the invention applies a stop-and-wait ARQ method on each one of the ARQ Channels A-F. Under this arrangement, the source 16 sequentially transmits data packets in multiplexing rounds, starting with the first ARQ



**FIG. 1**  
(PRIOR ART)

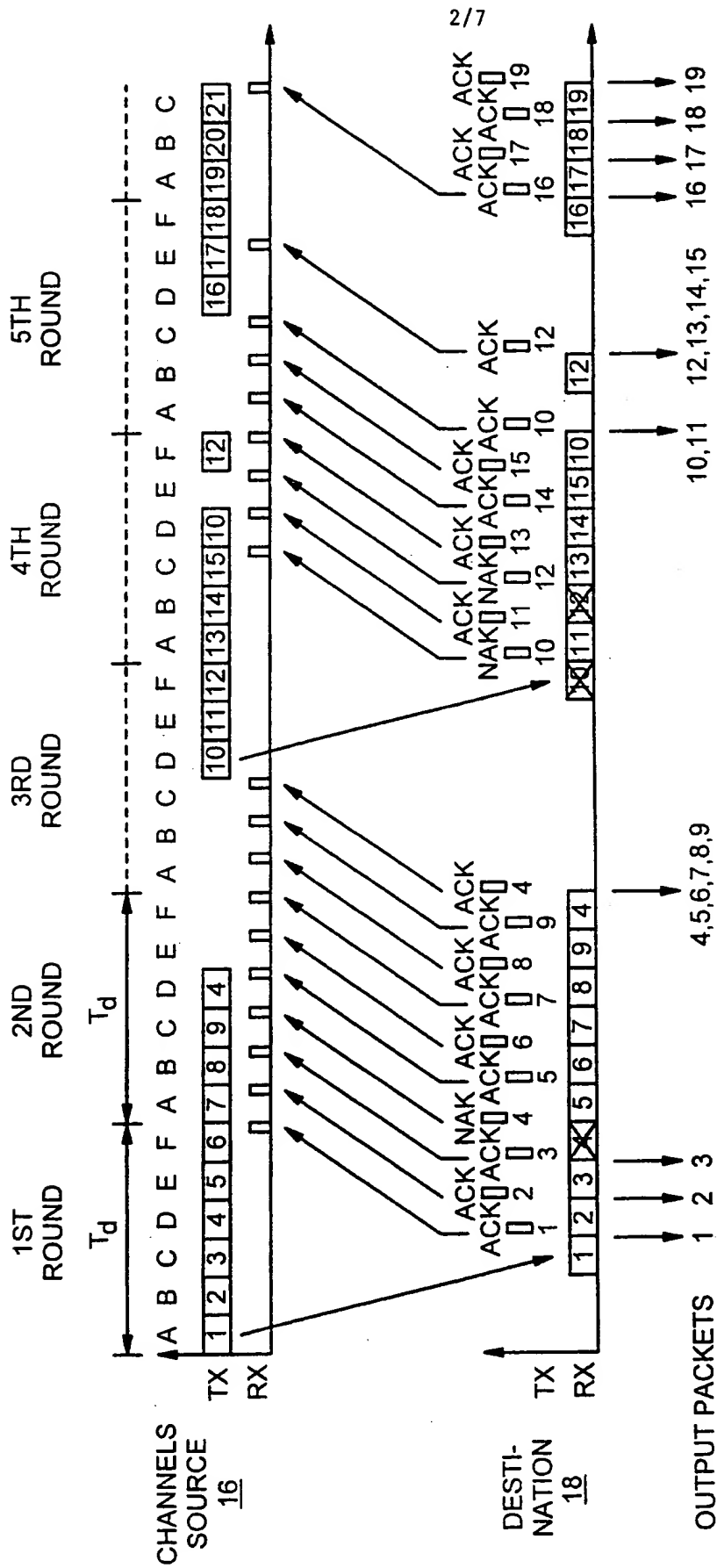


FIG. 2

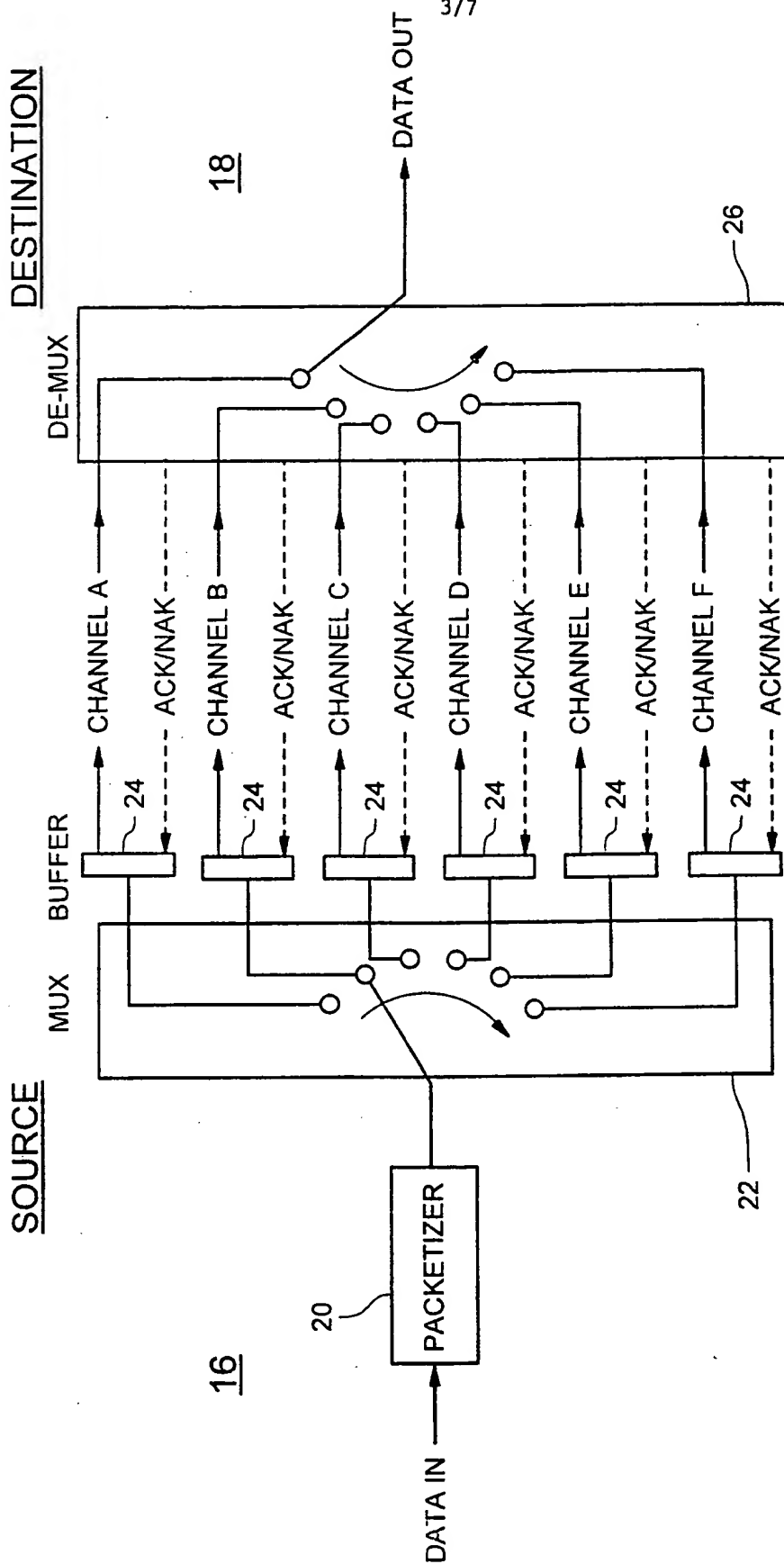


FIG. 3

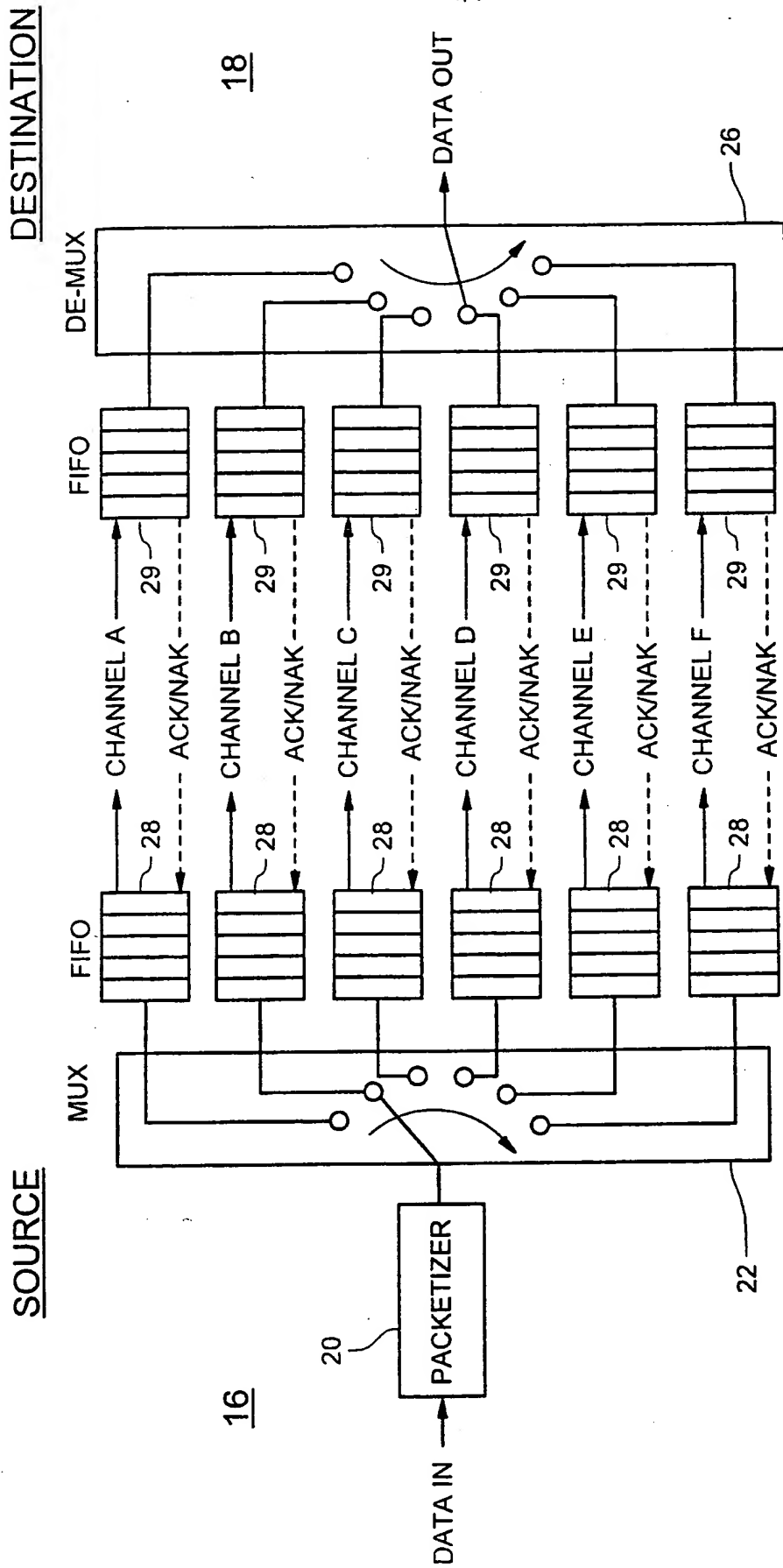


FIG. 4